Continuous-variable and hybrid quantum gates

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We provide several schemes to construct the continuous-variable SWAP gate and present a Hermitian
generalized many-body continuous controlled-\textsuperscript{a}-NOT gate. We introduce and study the hybrid
controlled-NOT gate and controlled-SWAP gate, and physical realizations of them are discussed
in trapped-ion systems. These continuous-variable and hybrid quantum gates may be used in the
corresponding continuous-variable and hybrid quantum computations.

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I. INTRODUCTION

The quantum computer [1,2] is a device which operates with quantum logic gates. It was shown that any quantum
computation can be built from a series of one-bit and two-bit quantum logic gates [3]. The fundamental
controlled-NOT (CN) [4] gate, widely discussed in the literature [5], is the two-qubit gate in which one qubit is
flipped conditioned on the state of another qubit. Mathematically the CN gate is defined as

\[ \text{CN}_{12} |i\rangle_1 |j\rangle_2 = |i\rangle_1 (|i\rangle + |j\rangle)_2, \]

(1)

where \(|i\rangle_j\rangle_2 \) \((i, j = 0, 1)\) are the basis states of the two
qubits, \(\oplus\) denotes addition modulo 2. The first (second) qubit is the control (target).

It is known that an unknown qubit state \(|\psi\rangle\) can be swapped with the qubit state \(|0\rangle\) using only two CN gates
[6], i.e.,

\[ \text{CN}_{21} \text{CN}_{12} |\psi\rangle_1 |0\rangle_2 = |0\rangle_1 |\psi\rangle_2. \]

(2)

In Ref. [7], the gate \(\text{CN}_{21} \text{CN}_{12}\) is also called double CN
gate. Using the CN gates one can construct a general
two-qubit SWAP gate as follows

\[ \text{SWAP}_{12} = \text{CN}_{12} \text{CN}_{21} \]

(3)

which makes the transformation

\[ \text{SWAP}_{12} |i\rangle_1 |j\rangle_2 = |j\rangle_1 |i\rangle_2. \]

(4)

The SWAP gate can be constructed in an alternative way
[8],

\[ \text{SWAP}_{12} = \frac{1}{2} (1 + \sigma_x \sigma_x + \sigma_y \sigma_y + \sigma_z \sigma_z), \]

(5)

where the operators \(\sigma_\alpha (\alpha = x, y, z)\) are the usual Pauli
operators of system \(i\). The remarkable properties of the
SWAP gate are described by Collins \textit{et al.} [7], Eisert \textit{et al.}
[9], and Chefles \textit{et al.} [10]. Both the CN gate and SWAP

gate are two-qubit gates. The one-qubit gates include
NOT gate which is expressed by the Pauli operator \(\sigma_x\)
and the Hadmard gate

\[ H = \frac{1}{\sqrt{2}} (\sigma_x + \sigma_z) \]

(6)

which makes the transformation

\[ H|0\rangle = \frac{1}{\sqrt{2}} (|0\rangle + |1\rangle), \]

(7a)

\[ H|1\rangle = \frac{1}{\sqrt{2}} (|0\rangle - |1\rangle). \]

(7b)

Both the NOT gate and the Hadmard gate are self-inverse, i.e., the square of them are the identity
operators.

For three qubits there are two types of gates, the
Toffoli [11] gate and Fredkin gate [12], which are also
called (controlled)\textsuperscript{2}-NOT gate and the controlled-SWAP
(CSWAP) gate, respectively. The CSWAP gate makes the following transformation

\[ \text{CSWAP}_{123} |i\rangle_1 |j\rangle_2 |0\rangle_3 = |i\rangle_1 |j\rangle_2 |0\rangle_3, \]

(8a)

\[ \text{CSWAP}_{123} |i\rangle_1 |j\rangle_2 |1\rangle_3 = |j\rangle_1 |i\rangle_2 |1\rangle_3, \]

(8b)

where the third qubit acts as the control. The quantum
gates described above act on discrete variables, the
qubits. In this paper we give the continuous-variable and
hybrid versions of the discrete quantum gates, which may
be used in the continuous-variable [13] and hybrid [14]
quantum computation. In the hybrid version of quantum
gates the discrete variable acts as the control and the
continuous variable as the target.

In Sec. II we begin with the introduction of the one-
body gates for continuous variables. We proceed in Sec.
III to study the two-body and many-body continuous-
variable gates and consider the CN gate, SWAP gate,
and controlled\textsuperscript{a}-NOT gate as well as the cloning gate.
Several methods are proposed to realize the SWAP gate. In Sec. III we introduce and study the hybrid quantum
gates, hybrid CN gates and CSWAP gates. We give two
schemes to realize the hybrid gates in trapped-ion sys-
tems. The conclusion is given in Sec. V.

II. ONE-BODY GATES FOR CONTINUOUS

A. NOT gate

The one-body continuous-variable NOT gate is the
parity operator, which is defined as

\[ H = \frac{1}{\sqrt{2}} (\sigma_x + \sigma_z) \]

(6)
\[ \text{NOT} = (-1)^a \sigma^a, \]  
where \( a \) and \( a^\dagger \) are bosonic annihilation and creation operators. It is easy to see that

\[
\text{NOT}|x\rangle = | -x\rangle, \\
\text{NOT}|p\rangle = | -p\rangle, \\
\text{NOT}^2 = 1, 
\]
where \( |x\rangle \) is the eigenstate of the position operator \( \hat{x} \), and \( |p\rangle \) is the eigenstate of the momentum operator \( \hat{p} \).

**III. TWO-BODY AND MANY BODY GATES FOR CONTINUOUS VARIABLES**

**A. CN gate**

The two-qubit CN gate has been extended to the case of continuous variables, the gates \( \text{CN}_{12}^+ \) [15] and \( \text{CN}_{12}^- \) [16], which are defined by

\[
\text{CN}_{12}^\pm |x_1\rangle_1 |y_2\rangle_2 = |x_1\rangle_1 |x \pm y\rangle_2, \\
\text{CN}_{12}^- = e^{-ix_1 \hat{p}_2}, \\
\text{CN}_{12}^+ = \text{NOT}_2 e^{ix_1 \hat{p}_2} e^{-ix_1 \hat{p}_2} \text{NOT}_2, 
\]
where the position operator of system \( i \) (\( i = 1, 2 \)) is denoted by \( \hat{x}_i \) and the momentum operator by \( \hat{p}_i \). In momentum space the CN gate can be defined as

\[
\text{CN}_{12}^\pm |p_1\rangle_1 |q_2\rangle_2 = |p_1\rangle_1 |p \pm q\rangle_2, \\
\text{CN}_{12}^- = e^{ix_2 \hat{p}_1}, \\
\text{CN}_{12}^+ = \text{NOT}_2 e^{-ix_2 \hat{p}_1} = e^{ix_2 \hat{p}_1} \text{NOT}_2. 
\]

The definitions of the CN gates are basis dependent. From Eqs. (14), (15), and (18), it is easy to check that both gates are unitary, the gate \( \text{CN}_{12}^- \) is not Hermitian and not self-inverse, while \( \text{CN}_{12}^+ \) is Hermitian and self-inverse.

The continuous version of the Hadamard gate is in fact the Fourier transformation and defined by [15]

\[
F(\sigma)|x\rangle = \frac{1}{\sigma \sqrt{\pi}} \int dy e^{2\pi ixy/\sigma^2} |y\rangle, 
\]

where \( \sigma \) is the scaled length. This is the transformation used to go from the position to the momentum basis if we set \( \sigma = \sqrt{2} \). The inverse \( F^\dagger(\sigma) \) is obtained by replacing \( i \) by \( -i \) giving the result that

\[
F(\sigma) F^\dagger(\sigma)|x\rangle = F^\dagger(\sigma) F(\sigma)|x\rangle = |x\rangle. 
\]

Note that the continuous-variable Hadamard gate is not self-inverse.

The CN gate for qubits has been used in various kinds of quantum information processing such as teleportation [17], dense coding [18], quantum state swapping [4], entangling quantum states [19] and Bell measurements [20]. It is natural to ask that if the continuous CN gates can perform some similar tasks like entangling and swapping quantum states. Let the continuous CN gates \( \text{CN}_{12}^\pm \) and the Hadamard gate \( F(\sqrt{2}) \) act on the state \( |z\rangle_1 |y\rangle_2 \). The resultant states are entangled states

\[
|\psi\rangle = \text{CN}_{12}^\pm F(\sqrt{2}) |z\rangle_1 |y\rangle_2 \\
= \frac{1}{\sqrt{2\pi}} \int dx e^{ixz} |x\rangle_1 |x \pm y\rangle_2. 
\]

It is interesting to see that the following equations

\[
(\hat{x}_1 - \hat{x}_2)|\psi\rangle = \mp y |\psi\rangle, \\
(\hat{p}_1 + \hat{p}_2)|\psi\rangle = z |\psi\rangle, 
\]

hold. That is to say, both the entangled states \( |\psi\rangle \) are the common eigenvectors of the position difference operator \( \hat{x}_1 - \hat{x}_2 \) and momentum sum operator \( \hat{p}_1 + \hat{p}_2 \). Further both the continuous CN gates can be used to construct \( N \)-party entangled state as follows

\[
\text{CN}_{12}^\pm \text{CN}_{13}^\pm \ldots \text{CN}_{1N}^\pm |p = 0\rangle_1 |x = 0\rangle_2 \\
= \frac{1}{\sqrt{2\pi}} \int dx_1 |x_1\rangle_2 \ldots |x\rangle_N. 
\]

This state is obtained by Braunstein [15] by a series of beam splitters. Here we provide an alternative way to obtain this state by using \( N \) CN gates. The \( N \)-party entangled state is an eigenstate with total momentum zero and relative positions zero.

**B. SWAP gate**

Having seen that both the continuous CN gates can entangle quantum states, then we ask if they can perform quantum state swapping by certain combinations of them. For continuous variables we have

\[
\text{CN}_{21}^- \text{CN}_{12}^+ |x\rangle_1 |y = 0\rangle_2 = |y = 0\rangle_1 |x\rangle_2. 
\]

From Eq.(3), one may guess that similar expression exists for continuous-variable SWAP gate. It is straightforward to check that

\[
\text{CN}_{12}^\pm \text{CN}_{21}^\pm |x\rangle_1 |y\rangle_2 = |2x + y\rangle_1 |3x + 2y\rangle_2, \\
\text{CN}_{12}^{-} \text{CN}_{21}^{-} |x\rangle_1 |y\rangle_2 = | -y\rangle_1 |-x\rangle_2. 
\]

Then the SWAP gate can be constructed as

\[
\text{SWAP}_{12} = \text{NOT}_1 \text{NOT}_2 \text{CN}_{12} \text{CN}_{21} \text{CN}_{12} \\
= \text{CN}_{12} \text{CN}_{21} \text{NOT}_1 \text{NOT}_2. 
\]
We see that one can not obtain the SWAP gate by only the gates $\text{CN}_{ij}^+(i \neq j)$, while one can use the gates $\text{CN}_{ij}^-$ to obtain it. Different from the situation of discrete variables, here the continuous-variable SWAP gate needs two NOT gates. In fact the gates $\text{CN}_{ij}^+(i \neq j)$ is not completely useless in the realization of the SWAP gate. Using both the gates $\text{CN}_{ij}^+$ and $\text{CN}_{ij}^-$, we have

$$\text{SWAP}_{12} = \text{NOT}_2\text{CN}_{12}^-\text{CN}_{21}^-\text{CN}_{12}^+ = e^{i\hat{p}_2\hat{p}_1}\text{NOT}_1 e^{i\hat{p}_2\hat{p}_1} e^{-i\hat{p}_2\hat{p}_1}. \quad (26)$$

Here we have used Eqs.(14) and (15). Then we can construct the SWAP gate using one-body gate and three two-body gates. The SWAP gate acting on momentum space can be constructed similarly.

Recalling that the two-qubit SWAP gate can be given in Eq.(5), we expect that the continuous SWAP gate be implemented in another way. Now we introduce the operator

$$B_{12} = e^{i\frac{\pi}{2}(\hat{p}_1\hat{p}_2 - \hat{2}_1\hat{2}_2)} \quad (27)$$

acting on the two continuous systems 1 and 2. It makes the transformation

$$B_{12} \left( \begin{array}{c} \hat{p}_1 \\ \hat{p}_2 \end{array} \right) B_{12}^{\dagger} = \left( \begin{array}{c} -\hat{p}_2 \\ \hat{p}_1 \end{array} \right), \quad (28)$$

from which we have

$$B_{12} |x\rangle |y\rangle_2 = |y\rangle |x\rangle_2. \quad (29)$$

Then the continuous-variable SWAP gate is immediately obtained as

$$\text{SWAP}_{12} = \text{NOT}_2 B_{12}. \quad (30)$$

From Eqs.(28) and (30), the swapping function of the SWAP gate can be compactly stated by

$$\text{SWAP}_{12} \left( \begin{array}{c} \hat{p}_1 \\ \hat{p}_2 \end{array} \right) \text{SWAP}_{12} = \left( \begin{array}{c} \hat{p}_2 \\ \hat{p}_1 \end{array} \right),$$

$$\text{SWAP}_{12} \left( \begin{array}{c} \hat{x}_1 \\ \hat{x}_2 \end{array} \right) \text{SWAP}_{12} = \left( \begin{array}{c} \hat{x}_2 \\ \hat{x}_1 \end{array} \right), \quad (31)$$

which may serve as alternative definitions.

Substituting $\tilde{x}_j = \frac{1}{\sqrt{2}}(a_j + a_j^\dagger)$, $\tilde{p}_j = \frac{1}{\sqrt{2}}(a_j - a_j^\dagger)$ to the Eq.(27), we can reexpress the operator $B_{12}$ in terms of the annihilation and creation operators and then rewrite the SWAP gate (30) as

$$\text{SWAP}_{12} = e^{i\pi a_1^\dagger a_2 e^{\frac{1}{2}a_1^\dagger a_2 - a_1^\dagger a_1}}. \quad (32)$$

Let the above SWAP gate act on the discrete Fock basis states, we obtain

$$\text{SWAP}_{12} |n\rangle_1 |m\rangle_2 = |m\rangle_1 |n\rangle_2, \quad (33)$$

where $|n\rangle_i$ denotes the Fock state of system $i$. Eq.(33) in fact gives the representation of the SWAP gate in the two-mode Fock space. We see that the SWAP gate is basis independent, while the CN gate is basis dependent.

As an end of this subsection, we see that there is a relation between the SWAP gates and the CN gates,

$$\text{SWAP}_{12}\text{CN}_{12}\text{SWAP}_{12} = \text{CN}_{21}. \quad (34)$$

The above equation shows that one can use the SWAP gate and CN gate $\text{CN}_{12}$ to realize another CN gate $\text{CN}_{21}$.

### C. Controlled*-NOT gate

We define a Hermitian continuous generalization of the discrete controlled*- NOT gate as

$$\text{CN}_{(12...N)+1}^\pm |x_1\rangle_1 |x_2\rangle_2 ... |x_N\rangle_N |x_{N+1}\rangle_{N+1} = |x_1\rangle_1 |x_2\rangle_2 ... |x_N\rangle_N - \sum_{n=1}^{N} |x_n\rangle_{N+1}, \quad (35)$$

$$\text{CN}_{(12...N)+1}^\pm = \text{NOT}_{N+1} e^{-i\hat{p}_{N+1} \sum_{n=1}^{N} \hat{x}_n}. \quad (36)$$

Similar gate can be defined in momentum space. Then the gate defined in this way is both unitary and Hermitian, and therefore self-inverse. For the case $N = 3$ and 2, the gate becomes the continuous-variable Toffoli and CN gate, respectively.

### D. 1 $\rightarrow$ 2 cloning gate

For qubits the CN gates $\text{CN}_{21}$ and $\text{CN}_{31}$ commute with each other, however for continuous variables, from Eq.(15), the following equation

$$[\text{CN}_{31}^-, \text{CN}_{21}^+] = e^{i(\tilde{x}_2 - \tilde{x}_3)}\hat{p}_1 - e^{i(\tilde{x}_3 - \tilde{x}_2)}\hat{p}_1 \quad (37)$$

holds for two Hermitian CN gates $\text{CN}_{21}^-$ and $\text{CN}_{31}^-$. That is to say, these two CN gates do not commute.

It is known that the 1 $\rightarrow$ 2 cloning gate is described by [21]

$$C = \text{CN}_{31}\text{CN}_{21}\text{CN}_{12}\text{CN}_{13} \quad (38)$$

in terms of four CN gates. To generalize directly to the continuous case of the above cloning gate, we obtain

$$C' = \text{CN}_{31}^\dagger\text{CN}_{21}^\dagger\text{CN}_{12}^\dagger \text{CN}_{13}^\dagger \quad (39)$$

Using Eqs.(15) and (37), we rewrite the gate $C'$ as

$$C' = e^{-i(\tilde{x}_3 - \tilde{x}_2)}\hat{p}_1 e^{-i\tilde{x}_1(\hat{p}_2 + \hat{p}_3)}\text{NOT}_2\text{NOT}_3. \quad (40)$$

which is just the continuous-variable 1 $\rightarrow$ 2 cloning gate up to the two NOT gates [22].
IV. HYBRID GATES

Now we introduce and study two kinds of hybrid quantum gates, the hybrid CN gate and CSWAP gate.

A. Hybrid CN gate

We define the hybrid CN gate as

\[
\begin{align*}
\text{CN}'_{12}(0)_{1}|x\rangle_{2} &= |0\rangle_{1}|x\rangle_{2}, \\
\text{CN}'_{12}(1)_{1}|x\rangle_{2} &= |1\rangle_{1}-|x\rangle_{2},
\end{align*}
\]

which can be realized in a trapped-ion system. In trapped-ion systems, one can have the following Hamiltonian experimentally [23, 24]

\[
H_1 = \lambda a^\dagger a P_1
\] (41)

where \(a\) and \(a^\dagger\) are bosonic annihilation and creation operators of the center-of-mass motion of the trapped ion, \(P_1 = |1\rangle_1\langle 1|\) is the projection operator, and \(\lambda\) is the effective coupling constant. It is easy to show that the evolution operator \(e^{-i\Delta t a^\dagger a P_1}\) at time \(t = \pi/\lambda\) gives directly the hybrid CN gate. One simple application of this gate is the generation of even and odd coherent states. Let the input state be \(\frac{1}{\sqrt{2}}(|0\rangle + |1\rangle_1)|\alpha\rangle_2\), where \(|\alpha\rangle_2\) is a bosonic coherent state. Then after the operation the output state will be \(\frac{1}{\sqrt{2}}(|0\rangle_1|\alpha\rangle_2 + |1\rangle_1-\alpha\rangle_2\). Now we measure the qubit on the state \(|\pm\rangle = \frac{1}{\sqrt{2}}(|0\rangle \pm |1\rangle)\), the continuous state will collapse into the even and odd coherent states, respectively.

B. Hybrid controlled-SWAP gate

A general controlled-SWAP gate is described by the following transformation

\[
\begin{align*}
|\Psi\rangle_1|\Phi\rangle_2|0\rangle_3 &\rightarrow |\Psi\rangle_2|\Phi\rangle_1|0\rangle_3, \\
|\Psi\rangle_1|\Phi\rangle_2|1\rangle_3 &\rightarrow |\Psi\rangle_2|\Phi\rangle_1|1\rangle_3, 
\end{align*}
\]

(42)

This gate has three inputs and the third is the control qubit. Let the input state of the CSWAP gate be \(\frac{1}{\sqrt{2}}|\Psi\rangle_1|\Phi\rangle_2(|0\rangle_3 + |1\rangle_3)\) and measure the output state. If we measure the qubit on the state \(|\pm\rangle = \frac{1}{\sqrt{2}}(|0\rangle \pm |1\rangle)\), we obtain exactly the symmetric and antisymmetric entangled states, \(|\Psi\rangle_1|\Phi\rangle_2 \pm |\Phi\rangle_2|\Psi\rangle_1\) up to normalization constants. This is actually a universal entangler [25]. So it is desirable to consider the CSWAP gate of the form (42) when then states \(|\Psi\rangle_1\) and \(|\Phi\rangle_2\) are continuous-variable states.

From the continuous-variable SWAP gate (26), the CSWAP gate is formally constructed as

\[
\text{CSWAP}'_{12(3)} = e^{ix_1\hat{p}_2P_3}e^{ix_2\hat{a}_1P_3}e^{ix_2\hat{a}_1P_3}e^{ix_1\hat{p}_2P_3}.
\] (43)

where \(P_3 = |1\rangle_3\langle 1|\) is the projection operator of the control system 3. There are three three-body interactions in the expression of the CSWAP gate. We will realize the CSWAP gate by two-body interactions.

First we see that the operators \(e^{i\frac{1}{2}x_i p}\), \(e^{\pm i\frac{3}{2}p}\) satisfy relation

\[
e^{ixp} = e^{ip\hat{p}}e^{ip\hat{p}}e^{-ip\hat{p}}e^{-ip\hat{p}}.
\] (44)

The above relation can be generalized as [26]

\[
e^{ixpsin\theta} = e^{i(\frac{x}{\theta})a e^{ip\hat{p}} e^{-i(\frac{x}{\theta})a e^{ip\hat{p}}} \times e^{i(\frac{1}{\theta})a e^{ip\hat{p}} e^{-i(\frac{1}{\theta})a e^{ip\hat{p}}}}.
\]

As the operator \(\hat{p}_1, \hat{x}_2\) and \(P_3\) commutes with each other, we replace \(x\) with \(\hat{x}_2\), \(\hat{x}\) with \(\hat{p}_1\), and \(\theta\) with \(\pi P_3 / 2\) in Eq. (45), respectively. Then we obtain

\[
e^{ip_1\hat{x}_2P_3} = e^{i\frac{3}{2}(1-P_3)a^\dagger a e^{ip_1}\hat{x} e^{-i\frac{3}{2}(1-P_3)a^\dagger a e^{ip_1}\hat{x}}\times e^{i\frac{1}{2}(1-P_3)a^\dagger a \hat{e} e^{-i\frac{1}{2}(1-P_3)a^\dagger a \hat{e}}}
\]

(46)

The above equation shows that we have written the three-body unitary operator \(e^{ip_i\hat{x}_2P_3}\) in terms of eight two-body operators. Therefore the CSWAP gate (43) can be written in terms of two-body operators.

From Eqs. (27) and (30), we write the CSWAP gate as the form

\[
\text{CSWAP}'_{12} = e^{i\pi a_1 a_2 P_3}e^{i\frac{3}{2}(1-P_3)\hat{x}_1\hat{p}_2\hat{x}_2P_3},
\] (47)

which also includes a three-body operator. Next we see how to realize this CSWAP gate in a trapped-ion system.

Gerry derived an effective Hamiltonian for two modes \(a\) and \(b\) as [27]

\[
H_2 = \chi(a_1^\dagger a_1 - a_2^\dagger a_2)P_3
\] (48)

in a trapped-ion system. The Hamiltonian \(H_2\) can be rewritten as

\[
H_2 = 2\chi J_z P_3,
\] (49)

where \(J_z = \frac{1}{2}(a_1^\dagger a_1 - a_2^\dagger a_2)\). The operators \(J_z, J_+ = a_1^\dagger a_2,\) and \(J_- = a_2^\dagger a_1\) form the \(\text{su}(2)\) Lie algebra. The unitary operator at time \(t = -\pi/(2\chi)\) corresponds to the Hamiltonian is given by

\[
U = e^{i\pi J_z P_3}
\] (50)

The unitary operator \(U\) can be transformed to \(U'\) as

\[
U' = e^{i\frac{3}{2}J_z}Ue^{-i\frac{3}{2}J_z} = e^{i\pi J_y P_3}e^{i\frac{3}{2}(\hat{x}_1\hat{p}_2-\hat{x}_2\hat{p}_1)P_3},
\] (51)

where \(J_x = (J_+ + J_-)/2\) and \(J_y = (J_+ - J_-) / (2i)\). From Eqs. (47) and (51), we write the CSWAP gate as

\[
\text{C-SWAP}'_{12} = e^{ia_1^\dagger a_2 P_3}e^{i\frac{3}{2}(1-P_3)(a_1^\dagger a_2 + a_2^\dagger a_1)}
\]

\[
\times e^{i\frac{3}{2}a_2^\dagger a_1 P_3}e^{-i\frac{3}{2}(1-P_3)a_1^\dagger a_2 P_3}e^{-i\frac{3}{2}(1-P_3)a_2^\dagger a_1 P_3},
\]

(52)

\[
\times e^{ia_1^\dagger a_2 P_3}e^{-i\frac{3}{2}(1-P_3)a_1^\dagger a_2 P_3}e^{-i\frac{3}{2}(1-P_3)a_2^\dagger a_1 P_3}.
\] (52)
Therefore we have given the form of CSWAP gate in terms of five two-body operators.

We have used two methods to express the three-body hybrid CSWAP gate in terms of two-body operators. In other words we provide two ways to realize the CSWAP gate.

V. CONCLUSION

In conclusion we have introduced and studied the continuous and hybrid versions of quantum gates. The continuous-variable gates include one-body (NOT, Hadamard), two-body (CN, double CN, SWAP) and many-body gates (controlled^n-NOT). Some relations between the CN, double CN and the SWAP gates are given. The hybrid quantum gates include the hybrid CN gate and the three-body controlled-SWAP gate. We proposed physical schemes to realize the hybrid gates in the trapped-ion systems. It is interesting to see that most of the quantum gates are not only unitary, but also Hermitian, and therefore self-inverse.

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