I. INTRODUCTION

Remote state preparation (RSP) is the preparation of a state at a remote location using entanglement and classical communication. In general, one may perform exactly faithful RSP, producing exactly the desired state, or asymptotically faithful RSP, where the fidelity approaches one as the number of states prepared approaches infinity. It is well known that it is not possible to perform exactly faithful RSP without entanglement. An infinite amount of classical information is required to exactly represent an arbitrary state, and therefore exact RSP would require an infinite amount of classical communication if there were no entangled resource. A method for exact RSP of a restricted ensemble of states is given in Ref. [2], and an alternative method for exact RSP of arbitrary states is given in Ref. [3]. Recently Ye et al. [7] showed that it is possible to perform exact RSP using any pure entangled state, provided the Schmidt number is equal to the system dimension. However, the proof given in Ref. [7] does not give a complete technique for performing this remote state preparation.

Here we give an explicit technique that is based upon an approximate technique without entanglement, and quantify how much classical communication is required for this scheme. In Sec. III we describe this scheme, then we discuss the classical communication required in Sec. III. We show that the initial entangled state must have maximal Schmidt number in Sec. IV and conclude in Sec. V.

II. EXPLICIT SCHEME

As in Ref. [7], the initial state is an entangled state of the form

$$|\alpha\rangle = \sum_{k=0}^{d-1} \alpha_k |k\rangle |k\rangle,$$

where the $\alpha_k$ are nonzero real numbers, and each subsystem is of dimension $d$. Any entangled state with maximal Schmidt number may be brought to this form via local operations. The state we wish to prepare is

$$|\beta\rangle = \sum_{k=0}^{d-1} \beta_k |k\rangle,$$

where the $\beta_k$ may be complex.

To explain this remote state preparation scheme, we first explain a simple approximate scheme that one would use if no entangled resource state were available. In this case, one would communicate an approximation of the coefficients $\beta_k$, and prepare a state based on those coefficients. To approximate $\beta_k$, note that the real and imaginary parts of $\beta_k$ will be numbers in the interval $[-1,1]$. We can approximate $\beta_k$ by dividing the interval $[-1,1]$ into $D$ subintervals

$$[-1,2/D-1), [2/D-1,4/D-1), \cdots, [1-2/D, 1].$$

We then denote the numbers of the subintervals that the real and imaginary parts of $\beta_k$ lie in as $n^r_k$ and $n^i_k$. That is,

$$n^r_k = \min\{D, |D(\text{Re}\beta_k + 1)/2| + 1\},$$

$$n^i_k = \min\{D, |D(\text{Im}\beta_k + 1)/2| + 1\}.$$  

The min takes account of the fact that the last subinterval is closed, so 1 lies in subinterval $D$. We may then approximate the real and imaginary parts of $\beta_k$ as

$$\text{Re}\beta_k \approx (2n^r_k - 1)/D - 1, \quad \text{Im}\beta_k \approx (2n^i_k - 1)/D - 1.$$  

The error in this approximation will be no more than $1/D$.

We may define a state corresponding to this approximation by

$$|\tilde{\beta}\rangle = \sum_{k=0}^{d-1} \{(2n^r_k - 1)/D - 1 + i(2n^i_k - 1)/D - 1\} |k\rangle.$$  

This state will satisfy

$$|||\beta\rangle - |\tilde{\beta}\rangle|| \leq \frac{\sqrt{2d}}{D}.$$
However, the state $|\beta'\rangle$ is not necessarily normalized; the state that is prepared will be the corresponding normalized state, $|\beta\rangle$. This state may be a slightly poorer approximation, but will still satisfy (see Appendix A)

$$|\langle\beta|\beta'\rangle|^2 \geq 1 - \frac{2d}{D^2}. \tag{8}$$

Without an entangled state, one would communicate the $2d$ numbers $n^k_1$ and $n^k_2$ using $2d\log D$ bits. Here we use the convention that log indicates logarithms base 2. We also use the convention that the number of “bits” is the logarithm base 2 of the number of messages, and need not be an integer. The preparer would initialize the system in the state $|0\rangle$, then apply a unitary operation $U$ such that the final state is $|\beta\rangle$.

In the case where an entangled state is available, one may initialize the system in an alternative state $|\psi\rangle$ that is close to $|0\rangle$, such that the operation $U$ takes the system to the exact state $|\beta\rangle$. We express the required initial state $|\psi\rangle$ as

$$|\psi\rangle = \sum_{k=0}^{d-1} \psi_k e^{\phi_k} |k\rangle. \tag{9}$$

In order to prepare this state, we first apply an entanglement transformation scheme to transform the entangled state $|A\rangle$ to a second state

$$|\Psi\rangle = \sum_{k=0}^{d-1} \psi_k |k\rangle |k\rangle. \tag{10}$$

The communication that is required depends on the entanglement transformation method that is used. There are a number of different methods of performing entanglement transformations [3, 4, 10], but there is the problem that most of these methods require local operations in subsystem 2 that are dependent on the state to be prepared.

It is possible to use the entanglement transformation scheme in Ref. [3], though this method requires communication of $\log d^2$ bits to communicate the permutation used. Via Caratheodory’s theorem one may restrict the number of possible permutations to $d^2 - 2d + 2$, indicating that the communication required is approximately $2 \log d$. However, the set of $d^2 - 2d + 2$ permutations is dependent on the state to be prepared, so it is still necessary to communicate $\log d^2$ bits.

Here we describe a straightforward method of determining a set of permutations that is independent of the state to be prepared. In general, in order to perform the entanglement transformation, it is necessary that $\alpha^2 < \psi^2$. Here we apply the slightly stronger condition that $\psi^2 > 1 - r^2(d - 1)$, where $r = \min \{\alpha_k\}$. This condition implies that the majorization relation holds (see Appendix B).

The entanglement transformation may be achieved via a two step process. First the state is transformed from $|A\rangle$ to the intermediate state

$$|\Phi\rangle = \sum_{k=0}^{d-1} \phi_k |k\rangle |k\rangle, \tag{11}$$

where $\phi_0 = \psi_0$ and $\phi_k = \sqrt{1 - \psi^2_k}/(d - 1)$ for $k > 0$. This entanglement transformation may be achieved using the measurement operators

$$A_k = \sqrt{p_k} \left( \sum_{l=1,l\neq k}^{d-1} \frac{\phi_l}{\alpha_l} |l\rangle \langle l| + \frac{\phi_k}{\alpha_0} |0\rangle \langle 0| + \frac{\phi_0}{\alpha_0} |k\rangle \langle k| \right), \tag{12}$$

for $k > 0$, and

$$A_0 = \sqrt{p_0} \left( \sum_{l=1}^{d-1} \frac{\phi_l}{\alpha_l} |l\rangle \langle l| + \frac{\phi_0}{\alpha_0} |0\rangle \langle 0| \right). \tag{13}$$

The probabilities $p_k = (\alpha^2_k - \psi^2_k)/(\phi^2_0 - \phi^2_k)$ for $k > 0$ and $p_0 = 1 - \sum_{k>0} p_k$. On obtaining measurement result $k$, if $k > 0$ it is necessary to swap states $|0\rangle$ and $|k\rangle$. The total number of measurement results is $d$, so the communication required is $\log d$.

This entanglement transformation is followed by an entanglement transformation to take the state from $|\Phi\rangle$ to $|\Psi\rangle$. In this case the measurement operators required are

$$B_k = \frac{1}{\sqrt{d-1}} \left( |0\rangle \langle 0| + \sum_{l=1}^{d-1} \frac{\psi_{l\oplus k}}{\phi_l} |l\rangle \langle l| \right), \tag{14}$$

where $k > 0$ (there is no measurement operator for $k = 0$). The notation $\oplus$ is used to indicate addition modulo $d - 1$ but excluding $0$ (i.e. $1 + [l + k - 1] \mod (d - 1)$). On obtaining measurement result $k$, it is necessary to perform a cyclic permutation of the states $|1\rangle$ to $|d - 1\rangle$. The total number of possible measurement results is $d - 1$, so the communication required is $\log (d^2 - d)$. Thus this method allows one to transform $|A\rangle$ to $|\Psi\rangle$ with communication of only $\log (d^2 - d)$.

One may then use the method applied in the proof of Theorem 1 of Ref. [2] to obtain the state $|\psi\rangle$. That is, one may apply the projection operators

$$P_k = \frac{1}{d} |\chi_k\rangle \langle \chi_k| \tag{15}$$

where

$$|\chi_k\rangle = \sum_l e^{i[(2\pi/d)kl - \varphi_l]} |l\rangle. \tag{16}$$

Upon obtaining measurement result $k$ one performs the local operation

$$C_k = \sum_l e^{i[(2\pi/d)kl]} |l\rangle \langle l|. \tag{17}$$

This step requires an additional $\log d$ bits of classical communication.
The final step is to perform the local operation in subsystem 1 to take the state from $|\psi\rangle$ to $|\beta\rangle$. Communication of the numbers $n_k^c$ and $n_k^c$ that specify this operation requires communication of $2d \log D$. To determine the value of $D$ necessary, note that we have required $\psi_0^2 \geq 1 - r^2(d - 1)$ in order to perform the entanglement transformation. As $\psi_0^2 = |\langle 0|\psi\rangle|^2 = |\langle \beta|\beta\rangle|^2$, $\psi_0^2$ is equal to the fidelity between the state to be prepared, $|\beta\rangle$, and the approximate state $|\beta\rangle$. From Eq. (18), the condition $\psi_0^2 \geq 1 - r^2(d - 1)$ will be satisfied for

$$D = \left\lceil \sqrt{\frac{2d}{r^2(d - 1)}} \right\rceil. \quad (18)$$

To summarize, the RSP scheme with entanglement is a three step process:

**Step 1:** Transform $|\Lambda\rangle$ to $|\Psi\rangle$ using the measurement operators $[12]$, $[13]$, and $[14]$. The communication required is $\log(d^2 - d)$.

**Step 2:** Apply the method given in the proof of Theorem 1 of Ref. [7] to prepare the unentangled state $|\psi\rangle$. This step requires $\log d$ bits of communication.

**Step 3:** Perform the unitary operation $U$ to transform $|\psi\rangle$ to $|\beta\rangle$. This step requires communication of the numbers $n_k^c$ and $n_k^c$ to determine the operation $U$, and therefore requires communication of $2d \log D$ bits.

**III. CLASSICAL COMMUNICATION REQUIRED**

The total classical communication for this scheme is approximately $3 \log d + 2d \log D$. The classical communication required for this scheme is least when the entangled state used is close to a maximally entangled state. The amount of classical communication required goes to infinity as the entanglement approaches zero; there is therefore a tradeoff, just as in the asymptotic schemes considered by Refs. [3], [4].

The classical communication required is shown in Fig. 1 for the case of a qubit. Comparing with the figure given in Refs. [3], [4], we can see that the classical communication is significantly larger than for asymptotically faithful RSP. In contrast to the asymptotic case, it is also possible for the classical communication to approach infinity even if the entanglement is not approaching zero. This is possible because one of the Schmidt coefficients can become arbitrarily small even if the entanglement does not.

One question that naturally arises is whether it is possible to perform this RSP scheme with less classical communication. The total classical communication required for steps 1 and 2 only scales as $\log d$. This communication is already small, and it is unlikely that it can be improved upon. However, the communication for the final step is $2d \log D$, which is much larger.

One may slightly reduce the communication required for step 3 by noting that the global phase is arbitrary, so we may take $\beta_0$ to be real. Then it is only necessary to approximate $2d - 1$ numbers, and we obtain the fidelity

$$|\langle \beta|\beta\rangle|^2 \geq 1 - \frac{2d - 1}{D^2}. \quad (19)$$

Then the slightly lower value of $D$ may be taken

$$D = \left\lceil \sqrt{\frac{2d - 1}{r^2(d - 1)}} \right\rceil, \quad (20)$$

and the total communication for step 3 is $(2d - 1) \log D$. This only gives a slight reduction in the communication required; an example for qubit states is given in Fig. 1.

It is also possible to use a more efficient coding of the state. One method is to record the sign of the real and imaginary parts of $\beta_k$, then use $n_k^c$ and $n_k^c$ to approximate the absolute values of $\text{Re}\beta_k$ and $\text{Im}\beta_k$. For large $d$ most of the $n_k^c$ and $n_k^c$ will be small, so it is more efficient to record the number of digits in the binary representations of $n_k^c$ and $n_k^c$, as well as those digits. The total communication required is then no more than (see Appendix [C])

$$(2d - 1) \left[ -\log(r\sqrt{d - 1}) + \log[\log D'] + 2 \right], \quad (21)$$

where

$$D' = \left\lceil \frac{2d - 1}{4r^2(d - 1)} \right\rceil. \quad (22)$$

The first term is the communication required for the digits, and the second term is the communication required.
for the numbers of digits. The third term includes a correction for rounding, as well as the communication required for the signs.

In assessing the scaling of each of the terms with $d$ it is necessary to assume a scaling for $r$. It is not possible to take $r$ to be independent of $d$, because $r \leq 1/\sqrt{d}$. If $r \propto 1/\sqrt{d}$, the first term in Eq. (21) scales approximately linearly with $d$, whereas the second term scales as $d \log \log d$, and therefore is dominant for large $d$. However, this situation is unlikely, because it would mean that the communication required for the number of digits in $n_k^r$ and $n_k^\ell$ is less than that for the digits themselves. It is more realistic to assume that $r$ decreases more rapidly than $1/\sqrt{d}$ (for example as $1/d$); this is because, for larger dimension, it is more likely that one of the Schmidt coefficients is exceptionally small. Under this assumption, the first term is dominant, as would be expected.

It is possible to perform the coding more efficiently than this, although the proof is not constructive. In general, in order to approximate a state with fidelity $1 - \epsilon^2$, it is necessary to have a set of states $\mathcal{M} = \{|\varphi_k\rangle\}$ such that for any state $|\beta\rangle$, the fidelity between $|\beta\rangle$ and some element of $\mathcal{M}$ is at least $1 - \epsilon^2$. To approximate the state, it is necessary to communicate the index $k$ of a state that has fidelity at least $1 - \epsilon^2$ with $|\beta\rangle$. It was shown in Ref. [11] that the number of states in $\mathcal{M}$ need be no greater than $(2.5/\epsilon)^{2d}$; here we apply a similar method to improve upon this bound.

Consider a set $\mathcal{M}$ that satisfies the condition that $|\langle \varphi_k | \varphi_l \rangle|^2 < 1 - \epsilon^2$ for $k \neq l$. The largest set satisfying this condition, $\mathcal{M}_{\text{max}}$, must also satisfy the fidelity condition. This is because, if any state $|\beta\rangle$ satisfied $|\langle \varphi_k | \beta \rangle|^2 < 1 - \epsilon^2$ for all $k$, it could be added and thereby increase the size of the set. Because no two states in $\mathcal{M}_{\text{max}}$ have fidelity as large as $1 - \epsilon^2$, no state can have fidelity as large as $1 - (\epsilon/2)^2$ with more than one member of $\mathcal{M}_{\text{max}}$. Thus the regions of states with fidelity at least $1 - (\epsilon/2)^2$ with different elements of $\mathcal{M}_{\text{max}}$ can not intersect. One may therefore determine an upper limit on the number of states in $\mathcal{M}_{\text{max}}$ by dividing the volume of the region of normalized states by the volume of the region of states that has fidelity at least $1 - (\epsilon/2)^2$ with some state $|\varphi\rangle$.

The region of allowed states is the surface of a hypersphere, and has volume $2\pi^d/(d - 1)!$. From Appendix A the volume of a region with fidelity at least $1 - (\epsilon/2)^2$ is $2\pi^d(\epsilon/2)^{2d-2}/(d - 1)!$. Therefore the number of states in the set $\mathcal{M}_{\text{max}}$ is no larger than $(2/e)^{2d-2}$. In order to be able to perform the entanglement transformations, we require fidelity at least $1 - \epsilon^2$, where $\epsilon = r\sqrt{d-1}$. Therefore, the communication required for this non-constructive coding scheme is no more than

$$ (2d - 2) \log(2/r\sqrt{d - 1}). $$

We may place a lower bound on the communication required in a similar way. To do this, we divide the total volume of the region of normalized states by the volume of the region of states with fidelity at least $1 - \epsilon^2$ (with an arbitrary state). Clearly, if the number of states in $\mathcal{M}$ were less than this, then there would be at least some states that did not have fidelity at least $1 - \epsilon^2$ with any state in $\mathcal{M}$. The volume of normalized states is $2\pi^d/(d - 1)!$, whereas the region of states with fidelity at least $1 - \epsilon^2$ has volume $2\pi^d(\epsilon)^{2d-2}/(d - 1)!$. Thus the total number of states can be no less than $(1/e)^{2d-2}$.

Taking $\epsilon = r\sqrt{d-1}$, the classical communication can be no less than

$$ (2d - 2) \log(1/r\sqrt{d - 1}). $$

The communication required for the non-constructive method is close to this, as it is no more than $2d - 1$ bits larger. In addition, the lower bound (24) is similar to the first term in Eq. (21); therefore, provided the first term in (21) is dominant, the explicit method that we described earlier is close to optimal.

As the classical communication for the rest of the scheme is $\log[d^2(d-1)]$, for exact remote state preparation schemes of this type, the total communication used can not be less than

$$ \log[d^2(d-1)] + (2d - 2) \log(1/r\sqrt{d - 1}), $$

and there will be a scheme that uses communication of

$$ \log[d^2(d-1)] + (2d - 2) \log(2/r\sqrt{d - 1}). $$

These expressions are plotted for the case of $d = 2$ in Fig. 1. There is only a few bits difference between (20) and (24), and the explicit scheme given before requires communication that is greater than both (20) and (24).

It must be emphasised that the lower bound (24) is not a lower bound for arbitrary exact remote preparation schemes. One reason is that it was derived from the requirement that a state must be specified with fidelity $1 - r^2(d - 1)$. In order for it to be possible to apply the entanglement transformation from $|A\rangle$ to $|\Psi\rangle$, it is only necessary that $\alpha^2 < \psi^2$. The volume of states satisfying this condition will, in most cases, be larger, so it will be possible to specify the state with less communication (though more communication will be required for the state transformation). It is also possible that there may be some very different remote state preparation scheme that uses less communication.

### IV. Schmidt Number Required

It is possible to obtain stronger results for the Schmidt number of the entangled state. For the RSP scheme outlined above the Schmidt number of the entangled state used must be maximal. It is possible to prove that this is necessary for arbitrary exact RSP schemes as follows. First, note that the above exact RSP scheme is equivalent to a local measurement performed in subsystem $A$, followed by a unitary transformation applied in subsystem $B$ that is based on information communicated from subsystem $A$. 
This is not the most arbitrary RSP scheme possible. In general, one may add local ancillas, perform local unitary transformations, local general measurements, and two-way communication. The POVMs used in each subspace may depend on the results of previous measurements. Let the initial state be

\[ |A\rangle = \sum_{k=0}^{d'-1} \alpha_k |k\rangle|k\rangle, \]

where \( d' < d \). Because the local unitary transformations and measurement operators on subsystem \( A \) commute with those on subsystem \( B \), we may combine the operators on subsystem \( A \) into a single operator \( M_A(\beta, \phi) \). This operator may depend on the state to be prepared, \( |\beta\rangle \), as well as the results of measurements, \( \phi \). The vector \( \phi \) contains the results of measurements performed in both subsystems. We allow \( \phi \) to contain real numbers resulting from measurements in both subsystems (even though these results cannot be communicated with finite classical communication), as this does not make the RSP scheme less general. We also combine the operators on subsystem \( B \) into a single operator \( M_B(n, \phi) \). This operator also may depend on the results of measurements \( \phi \), as well as additional information \( n \) communicated from subsystem \( A \).

After performing operation \( M_A(\beta, \phi) \), the reduced density matrix in subsystem \( B \) is

\[ \rho \otimes \rho_{\text{anc}}, \]

where \( \rho_{\text{anc}} \) is the state of the ancilla for subsystem \( B \). As the ancilla for subsystem \( B \) is initially unentangled, it can not be modified in any way by \( M_A(\beta, \phi) \). In addition, although \( \rho \) will depend on \( M_A(\beta, \phi) \), it must still be orthogonal to \( |k\rangle \) for \( k > d' - 1 \). Without loss of generality, we assume that it is possible to prepare any state \( \rho \), provided it is orthogonal to \( |k\rangle \) for \( k > d' - 1 \). In order to obtain perfect RSP, we require

\[ |\beta\rangle \langle \beta| = \text{Tr}_{\text{anc}} \left[ M_B(n, \phi)(\rho \otimes \rho_{\text{anc}})M_B^\dagger(n, \phi) \right]. \]

If Eq. (27) holds for \( \rho \) and \( \rho_{\text{anc}} \), there must be pure states for which it holds. Therefore we may take these states to be \( |\psi\rangle \) and \( |\psi_{\text{anc}}\rangle \). Eq. (27) then becomes

\[ |\beta\rangle \otimes |\psi_{\text{anc}}\rangle = M_B(n, \phi)|\psi\rangle \otimes |\psi_{\text{anc}}\rangle , \]

where \( |\psi_{\text{anc}}\rangle \) is the final state of the ancilla.

In order to obtain \( |\beta\rangle \), for any given measurement results \( \phi \), one may adjust \( |\psi\rangle \) and the communicated information \( n \). An arbitrary pure \( d'-\text{dimensional state} \) \( |\beta\rangle \) is equivalent to a point on a \( 2d' - 1 \) dimensional hypersphere (one dimension may be omitted because we may take \( \beta_0 \) to be real). Because \( |\psi_{\text{anc}}\rangle \) is fixed, and \( |\psi\rangle \) is orthogonal to \( |k\rangle \) for \( k > d' - 1 \), the state \( |\psi\rangle \otimes |\psi_{\text{anc}}\rangle \) is equivalent to a point on a \( 2d' - 1 \) dimensional hypersphere. Since there is only a finite set of messages that may be communicated \( n \), the set of states obtained by varying \( n \) and \( |\psi\rangle \) can only correspond to a \( 2d' - 2 \) dimensional space, and cannot fill the \( 2d - 2 \) dimensional space corresponding to the set of states \( |\beta\rangle \).

Therefore, even if it is possible to prepare an arbitrary \( d'-\text{dimensional state} \) and perform one of a finite number of operations, it is not possible to prepare an arbitrary \( d\)-dimensional state. Thus it is not possible to exactly prepare an arbitrary \( d \)-dimensional state if the resource state has lower Schmidt number.

V. CONCLUSIONS

We have given an explicit scheme for performing exact RSP using an arbitrary entangled state with maximal Schmidt number and classical communication that is close to optimal for schemes of this type. The scheme is a three step process, involving an entanglement transformation, followed by a disentangling measurement and a final unitary operation to obtain the exact state.

This method improves on given in Ref. [7] in two main ways:

1. The communication required for the entanglement transformation is less than \( 2 \log d \), as compared to \( \log d! \) for Ref. [7].
2. We have given an explicit method for determining the final unitary operation.

The majority of the communication is required for the final unitary operation. The communication required for this step is slightly superlinear in \( d \), whereas the communication required for the first two steps is logarithmic in \( d \). This communication is close to optimal, provided the remote state preparation scheme is of this type; however, we have not eliminated the possibility that some more general remote state preparation scheme may require less communication.

This remote state preparation scheme also requires that the Schmidt number of the initial entangled state be maximal. We have proven that this is necessary even for an arbitrary remote state preparation scheme.

APPENDIX A: DISTANCE AND FIDELITY

Consider two states that satisfy

\[ |||\beta\rangle - |\beta'\rangle|| \leq \epsilon, \]

where \( |\beta'\rangle \) is not necessarily normalized. The state \( |\beta'\rangle \) may be expressed as \( |\beta'\rangle = a|\beta\rangle + b|\beta'\rangle \) where \( |\beta'\rangle \) is orthogonal to \( |\beta\rangle \). Then Eq. (A1) is equivalent to \( |1 - a|^2 + |b|^2 \leq \epsilon^2 \), which implies

\[ |a| \geq 1 - \sqrt{\epsilon^2 - |b|^2}, \]

(A2)
The right-hand side of this expression is minimized for $|\beta|^2 = \epsilon^2 - \epsilon^4$, giving

$$\frac{|b|^2}{|a|^2} \geq \frac{\epsilon^2}{1 - \epsilon^2}. \quad \text{(A4)}$$

In turn this implies

$$\frac{|a|^2}{|a|^2 + |b|^2} \geq 1 - \epsilon^2. \quad \text{(A5)}$$

If $|\beta'|$ is the normalized state corresponding to $|\tilde{\beta}'\rangle$, then

$$|\langle \beta|\beta'\rangle|^2 = \frac{|a|^2}{|a|^2 + |b|^2}. \quad \text{(A6)}$$

Therefore $|||\beta| - |\tilde{\beta}'\rangle|| \leq \epsilon$ implies that $|\langle \beta|\beta'\rangle|^2 \geq 1 - \epsilon^2$.

**APPENDIX B: MAJORIZATION AND FIDELITY**

In this appendix it is shown that $\vec{\alpha} \prec \vec{\psi}^2$ is satisfied if $\psi_0^2 \geq 1 - (d - 1)r^2$. The majorization condition $\vec{\alpha} \prec \vec{\psi}^2$ is equivalent to

$$\sum_{k=0}^{p} 1\psi_k^2 \geq \sum_{k=0}^{p} 1\alpha_k^2, \quad \text{(B1)}$$

where the down arrow indicates that the coefficients are sorted into descending order. To show this result, note that, because the $1\psi_k^2$ are in descending order,

$$\frac{1}{p} \sum_{k=1}^{p} 1\psi_k^2 \geq \frac{1}{d - p - 1} \sum_{k=p+1}^{d-1} 1\psi_k. \quad \text{(B2)}$$

Multiplying on both sides by $d - p - 1$ and adding $\sum_{k=1}^{d-1} 1\psi_k^2$ gives

$$\frac{d - 1}{p} \sum_{k=1}^{p} 1\psi_k^2 \geq (1 - 1\psi_0^2). \quad \text{(B3)}$$

In turn this gives

$$\sum_{k=0}^{p} 1\psi_k^2 \geq 1\psi_0^2 \frac{d - p - 1}{d - 1} + \frac{p}{d - 1}. \quad \text{(B4)}$$

The substituting the inequality $\psi_0^2 \geq 1 - (d - 1)r^2$ (and using $1\psi_0^2 \geq \psi_0^2$) gives

$$\sum_{k=0}^{p} 1\psi_k^2 \geq 1 - (d - p - 1)r^2. \quad \text{(B5)}$$

Because $\alpha_k^2 \geq r^2$, it is also the case that

$$1 - (d - p - 1)r^2 \geq \sum_{k=0}^{p} 1\alpha_k^2, \quad \text{(B6)}$$

thus giving

$$\sum_{k=0}^{p} 1\psi_k^2 \geq \sum_{k=0}^{p} 1\alpha_k^2. \quad \text{(B7)}$$

Hence the inequality $\psi_0^2 \geq 1 - (d - 1)r^2$ is sufficient to imply the majorization relation $\vec{\alpha}^2 \prec \vec{\psi}^2$.

**APPENDIX C: EFFICIENT CODING**

If the numbers $n_k^r$ and $n_k^i$ record the absolute values of the real and imaginary parts of $\beta_k$, and the interval $[0, 1]$ is divided into $D'$ subintervals, then the fidelity is

$$\left| \langle \beta|\beta'\rangle \right|^2 \geq 1 - \frac{2d - 1}{4D'^2}. \quad \text{(C1)}$$

The number of subintervals should therefore be taken to be

$$D' = \left[ \sqrt{\frac{2d - 1}{4r^2(d - 1)}} \right]. \quad \text{(C2)}$$

The number of bits required to encode the length of the bit-strings representing each of the numbers $n_k^r$ and $n_k^i$ is $\log[\log D']$. In addition $\beta_0$ is taken to be real, and we require $2d - 1$ bits to record the signs of the real and imaginary parts of $\beta_k$. The total communication is therefore

$$(2d - 1) \log[\log D'] + \log(n_0^r - 1)$$

$$+ \sum_{k=1}^{d-1} \left[ \log(n_k^r - 1) + \log(n_k^i - 1) \right] + (2d - 1)$$

$$\leq (2d - 1) \log[\log D'] + \log(D\beta_0)$$

$$+ \sum_{k=1}^{d-1} \left[ \log(D\text{Re}\beta_k) + \log(D\text{Im}\beta_k) \right] + (2d - 1)$$

$$\leq (2d - 1) \left[ -\log(r\sqrt{d - 1}) + \log[\log D'] + 2 \right]. \quad \text{(C3)}$$

**APPENDIX D: VOLUME OF REGION OF STATES**

Here we consider the problem of determining the volume of the region of states $|\beta\rangle$ for a given $|\phi\rangle$ that satisfy $|\langle \phi|\beta\rangle|^2 \geq 1 - \epsilon^2$. To do this, we write the state $|\beta\rangle$ in the form

$$|\beta\rangle = e^{i\phi} \cos \theta|\phi\rangle + \sin \theta|\phi^+\rangle, \quad \text{(D1)}$$

$e^{i\phi} \cos \theta$ and $\sin \theta$ are uniformly distributed on $[0, 2\pi)$ and $[0, \pi/2)$ respectively.
where $|\varphi^\perp\rangle$ is a state perpendicular to $|\varphi\rangle$. Every state may be represented in this way when the ranges of $\phi$ and $\theta$ are $[-\pi, \pi]$ and $[0, \pi/2]$, respectively. The condition $|(\varphi|\beta\rangle|^2 \geq 1 - \epsilon^2$ implies that $|\sin \theta| \leq \epsilon$. The volume of states is given by

$$V = \int_0^{\arcsin \epsilon} d\theta \int_{-\pi}^{\pi} d\phi \cos \theta S_{2d-2}(\sin \theta),$$

(D2)

where

$$S_n(r) = \frac{2\pi^{n/2}}{\Gamma(n/2)} r^{n-1}$$

(D3)

is the surface area of a hypersphere. Integrating over $\phi$ and using (D3) gives

$$V = \frac{4\pi^d}{(d-2)!} \int_0^{\arcsin \epsilon} \cos \theta \sin^{2d-3} \theta d\theta$$

$$= \frac{4\pi^d}{(d-2)!} \left[ \frac{\sin^{2d-2} \theta}{2d-2} \right]_0^{\arcsin \epsilon}$$

$$= \frac{2\pi^d}{(d-1)!} \epsilon^{2d-2}.$$  

(D4)

Note that using $\epsilon = 1$ recovers the formula for the surface area of a $2d$ dimensional hypersphere, which gives the total volume of normalized states.

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[11] If $|\beta\rangle$ had fidelity as large as $1 - (\epsilon/2)^2$ with both $|\varphi_k\rangle$ and $|\varphi_l\rangle$, by the chain rule for fidelities, the fidelity between $|\varphi_k\rangle$ and $|\varphi_l\rangle$ would have to be at least $1 - \epsilon^2$. 